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**In the Claims:**

Claim 1. (Amended) A method for identification, comprising the steps of:

(a) generating system parameters  $G_1, G_2, P$  and  $\hat{e}$  and storing the system parameters in a memory by a system administrator, wherein  $G_1$  and  $G_2$  are cyclic groups of order  $m$ ,  $P$  is a generator on the cyclic group  $G_1$ ,  $\hat{e}$  is a bilinear map defined as

$$\hat{e}: G_1 \times G_1 \mapsto G_2;$$

(b) generating a private key  $\langle a, b, c \rangle$  and a public key  $v$  and storing the public key  $v$  in the memory by a prover or the system administrator, wherein  $a, b$  and  $c$  are randomly chosen in  $Z_m^*$  where  $Z_m^*$  is a multiplicative group of order  $m$ ;

(c) generating random numbers  $r_1, r_2, r_3 \in Z_m^*$  for obtaining an evidence  $(x, Q)$  and sending the evidence  $(x, Q)$  to a verifier by the prover;

(d) receiving the evidence  $(x, Q)$ , selecting a randomly selected number  $\omega \in Z_m^*$  to obtain a query  $R$ , storing the evidence  $(x, Q)$  and the randomly selected number  $\omega$  in the memory and sending the query  $R$  to the prover by the verifier;

(e) receiving the query  $R$ , computing a temporary value  $S$  to obtain a response  $Y$  and sending the response  $Y$  to the verifier by the prover; and

(f) determining a legitimacy of the prover by employing the system parameters  $G_1, G_2, P$  and  $\hat{e}$ , the public key  $v$ , the evidence  $(x, Q)$  and the randomly selected number  $\omega$  by the verifier.

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Claim 2. (Original)      The method of claim 1, wherein, in the step (b), the public key

$$v = \hat{e}(P, P)^{abc}$$

v is obtained by

Claim 3. (Original)      The method of claim 2, wherein, in the step (c), the evidence

$$x = \hat{e}(P, P)^{r_1 r_2 r_3}$$

(x, Q) includes a first evidence value and a second

$$Q = r_1 r_2 r_3 P$$

evidence value

Claim 4. (Original)      The method of claim 3, wherein, in the step (d), the query R is

$$R = \omega P$$

obtained by

Claim 5. (Original)      The method of claim 4, wherein, in the step (e), the temporary

$$S = r_1 r_2 r_3 R$$

value S is obtained by and the response Y is obtained by

$$Y = abcP + (a+b+c)S$$

Claim 6. (Original)      The method of claim 5, wherein the verifier determines the legitimacy of the prover by verifying

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$$\begin{aligned}\hat{e}(Y, P) &= \hat{e}(abcP + (a+b+c)S, P) \\&= \hat{e}(abcP + (a+b+c)r_1r_2r_3R, P) \\&= \hat{e}(abcP + (a+b+c)r_1r_2r_3\omega P, P) \\&= \hat{e}((abc + (a+b+c)r_1r_2r_3\omega)P, P) \\&= \hat{e}(P, P)^{abc + (a+b+c)r_1r_2r_3\omega} \\&= \hat{e}(P, P)^{abc} \cdot \hat{e}(P, P)^{(a+b+c)r_1r_2r_3\omega} \\&= \hat{e}(P, P)^{abc} \cdot \hat{e}(P, r_1r_2r_3P)^{(a+b+c)\omega} \\&= \hat{e}(P, P)^{abc} \cdot \hat{e}(P, Q)^{(a+b+c)\omega} \\&= \hat{e}(P, P)^{abc} \cdot \hat{e}((a+b+c), PQ)^{\omega} \\&= \hat{e}(P, P)^{abc} \cdot \hat{e}(aP+bP+cP, Q)^{\omega} \\&= v \cdot \hat{e}(aP+bP+cP, Q)^{\omega}\end{aligned}$$

Claim 7. (Amended)

A method for identification, comprising the steps of:

(a) generating system parameters  $G_1, G_2, P$  and  $\hat{e}$  and storing the system parameters in a memory by a system administrator, wherein  $G_1$  and  $G_2$  are cyclic groups of order  $m$ ,  $P$  is a generator on the cyclic group  $G_1$ ,  $\hat{e}$  is a bilinear map defined as

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$\hat{e}: G_1 \times G_1 \mapsto G_2$

(b) generating a private key  $\langle a_1, a_2, \dots, a_n \rangle$  and a public key  $v$  and storing the public key  $v$  in the memory by a prover or the system administrator, wherein  $a_1, a_2, \dots, a_n$  are randomly chosen in  $Z_m^*$  where  $Z_m^*$  is a multiplicative group of order  $m$ ;

(c) generating random numbers  $r_1, r_2, \dots, r_n \in Z_m^*$  for obtaining an evidence  $(x, Q)$  and sending the evidence  $(x, Q)$  to a verifier by the prover;

(d) receiving the evidence  $(x, Q)$ , selecting a randomly selected number  $\omega \in Z_m^*$  to obtain a query  $R$ , storing the evidence  $(x, Q)$  and the randomly selected number  $\omega$  in the memory and sending the query  $R$  to the prover by the verifier;

(e) receiving the query  $R$ , computing a temporary value  $S$  to obtain a response  $Y$  and sending the response  $Y$  to the verifier by the prover; and

(f) determining a legitimacy of the prover by employing the system parameters  $G_1, G_2, P$  and  $\hat{e}$ , the public key  $v$ , the evidence  $(x, Q)$  and the randomly selected number  $\omega$  by the verifier.

Claim 8. (Original)      The method of claim 7, wherein, in the step (b), the public key  $v$  is obtained by  $v = \hat{e}(P, P)^{a_1 a_2 \dots a_n}$ .

Claim 9. (Original)      The method of claim 8, wherein, in the step (c), the evidence  $(x, Q)$  includes a first evidence value  $v = \hat{e}(P, P)^{r_1 r_2 \dots r_n}$  and a second evidence value  $Q = r_1 r_2 \dots r_n P$ .

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Claim 10. (Original)      The method of claim 9, wherein, in the step (d), the query  $R$  is

$$R =_0 P$$

obtained by

Claim 11. (Original)      The method of claim 10, wherein, in the step (e), the temporary value  $S$  is obtained by  $S = r_1r_2\dots r_n R$  and the response  $Y$  is obtained by  $Y = a_1a_2\dots a_n P + (a_1+a_2+\dots+a_n)S$

Claim 12. (Original)      The method of claim 11, wherein the verifier determines the legitimacy of the prover by verifying

$$\begin{aligned} \hat{e}(Y, P) &= \hat{e}(a_1a_2\dots a_n P + (a_1+a_2+\dots+a_n)S, P) \\ &= \hat{e}(a_1a_2\dots a_n P + (a_1+a_2+\dots+a_n)r_1r_2\dots r_n R, P) \\ &= \hat{e}(a_1a_2\dots a_n P + (a_1+a_2+\dots+a_n)r_1r_2\dots r_n \omega P, P) \\ &= \hat{e}((a_1a_2\dots a_n + (a_1+a_2+\dots+a_n)r_1r_2\dots r_n \omega)P, P) \\ &= \hat{e}(P, P)^{a_1a_2\dots a_n + (a_1+a_2+\dots+a_n)r_1r_2\dots r_n \omega} \\ &= \hat{e}(P, P)^{a_1a_2\dots a_n} \cdot \hat{e}(P, P)^{(a_1+a_2+\dots+a_n)r_1r_2\dots r_n \omega} \\ &= \hat{e}(P, P)^{a_1a_2\dots a_n} \cdot \hat{e}(P, r_1r_2\dots r_n P)^{(a_1+a_2+\dots+a_n)\omega} \\ &= \hat{e}(P, P)^{a_1a_2\dots a_n} \cdot \hat{e}(P, Q)^{(a_1+a_2+\dots+a_n)\omega} \\ &= \hat{e}(P, P)^{a_1a_2\dots a_n} \cdot \hat{e}((a_1+a_2+\dots+a_n), PQ)^\omega \\ &= \hat{e}(P, P)^{a_1a_2\dots a_n} \cdot \hat{e}(a_1P+a_2P+\dots+a_nP, Q)^\omega \\ &= v \cdot \hat{e}(a_1P+a_2P+\dots+a_nP, Q)^\omega. \end{aligned}$$